

GRAPHICAL REPRESENTATIONS OF TOPOLOGIES ON A FINITE SET

by

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DEDICATION

I dedicate this thesis to my best friend, Brett Fuller, for letting me go, for letting me create, and for never expecting me to say the right thing. No one else has been worth the return trip. No one else inspires me more. And no one else has made me wish so hard that I could find the words.

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1. INTRODUCTION

As an undergraduate at Texas State University-San Marcos in an introductory topology course, Dr. Sukhjit Singh presented an exercise to the class of finding all the topologies on a three-element set.

Definition 1. A **topology** on a set X is a collection T of subsets of X having the following properties:

1. \emptyset and X are in T .
2. The union of elements of any subcollection of T is in T .
3. The intersection of the elements of any finite subcollection of T is in T . [8]

Most of my classmates tried to use brute force to tackle the problem. The brute force method would require the student to look at

$P(P(X)) = P(\{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\})$, which has a total of 256 elements to check for the properties of a topology.

Rather than perform such a time-consuming process, I drew a graph, G . I found it necessary to make vertices of two types. The first type was the single element subsets of $P(X) - \{\emptyset, X\} = \{\{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}\}$. So

$$V_1(G) = \{\{\{a\}\}, \{\{b\}\}, \{\{c\}\}, \{\{a, b\}\}, \{\{a, c\}\}, \{\{b, c\}\}\}.$$

The second type was the two element subsets of $P(X) - \{\emptyset, X\}$ where both elements have the same cardinality. So we have that $V_2(G)$ is the set:

$$\{\{\{a\}, \{b\}\}, \{\{a\}, \{c\}\}, \{\{b\}, \{c\}\}, \{\{a, b\}, \{a, c\}\}, \{\{a, b\}, \{b, c\}\}, \{\{a, c\}, \{b, c\}\}\}.$$

As this notation is cumbersome, we will often make use of the following representations:

$$\begin{aligned} \{\{a\}\} & \text{ is represented by } a \\ \{\{a,b\}\} & \text{ is represented by } ab \\ \{\{a\}, \{b\}\} & \text{ is represented by } a : b \\ \{\{a,b\}, \{a,c\}, \{b,c\}\} & \text{ is represented by } ab : ac : bc. \end{aligned}$$

Thus

$$\begin{aligned} V(G) &= V_1(G) \cup V_2(G) \\ &= \{a, b, c, ab, ac, bc, a : b, a : c, b : c, ab : ac, ab : bc, ac : bc\}. \end{aligned}$$

The edge set of G , $E(G)$, is define by the rule $v_1 \leftrightarrow v_2$ for some $v_1, v_2 \in V(G)$ if $v_1 \cup v_2 \cup \{\emptyset, X\}$ is a topology on X . For example, we have an edge from a to bc because $a \cup bc \cup \{\emptyset, X\} = \{\{a\}\} \cup \{\{b,c\}\} \cup \{\emptyset, X\} = \{\emptyset, \{a\}, \{b,c\}, \{a,b,c\}\}$. My original sketch is seen in Figure 1.

Now consider the topologies on X . The only topology with 2 open sets is the indiscrete topology $\{\emptyset, X\}$. The topologies with three open sets are all accounted for by $\{v \cup \{\emptyset, X\} \mid v \in V_1(G)\}$. Topologies with four open sets are represented on the graph by the edges between vertices in $V_1(G)$. The topologies with five and six open sets are then the edges drawn between vertices in $V_1(G)$ and $V_2(G)$ or edges drawn between two vertices of $V_2(G)$, respectively. The inclusion of any more sets beyond six would force the topology to be the discrete topology.

As this was an exercise to count the number of topologies on a three element set, I first counted the number of edges drawn, arriving at a total of 21 topologies. The elements of $V_1(G)$, when unioned with $\{\emptyset, X\}$, were also topologies in their own right, as mentioned above. Counting each of these, this added another six topologies for a total of 27. Only the

discrete and indiscrete topologies were not pictured in the graph. These two more topologies brought me to the final total of 29 topologies on a set of three elements.

The effectiveness of this construction for the purpose of counting topologies became an interesting question to me that I would continue to ponder throughout my undergraduate career: would the use of graphical representations be able to aid in the unanswered question of how to count topologies on a finite set? Even if not, what could be gleaned about the construction of topologies from the representative graphs? It was these questions which ultimately led me to a recursive definition for T_0 topologies on a finite set.

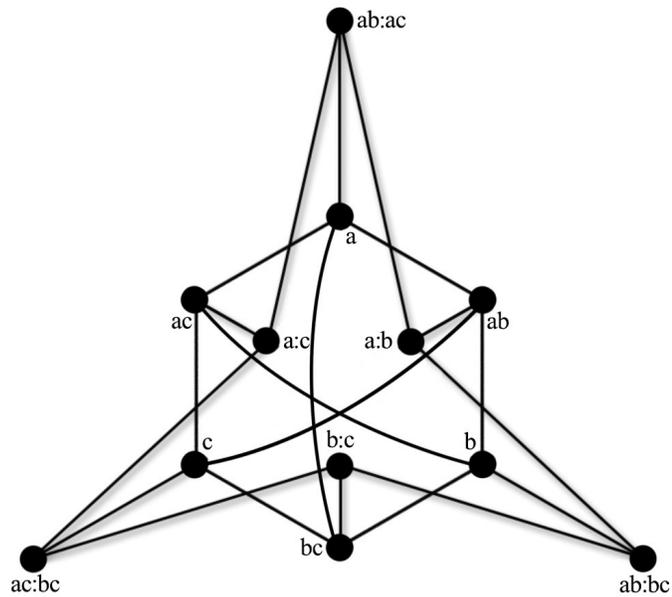


Figure 1: Original Graph of the Topologies on $X = \{a, b, c\}$

2. TERMINOLOGY AND NOTATIONAL CONVENTIONS

Throughout the course of this paper, we will make use of terminology and notation from the fields of graph theory and topology. When the construction is the work of the author, the definitions will be given in text. Presented here are standard definitions and notation considered background material for the following chapters.

We begin with an exploration of graphical representations. As described by West, a **graph** G is a triple consisting of a **vertex set** $V(G)$, an **edge set** $E(G)$, and a relation that associates with each edge two vertices (not necessarily distinct) called its **endpoints** [14]. Early representations used, such as the nucleus and shell graphs of Chapter 3, are undirected graphs. Using the elements of the **power set** of X , or the set of all possible subsets of X , we define a vertex set. For the relation which defines our edge set we turn to the definition of a topology, given in Definition 1.

This early work will focus largely on **distinct**, or **labeled**, **topologies**, which is a collection of all of the unique permutations of the elements of X into the open sets of a topology T . For example, when considering labeled topologies, the collections $\{\emptyset, \{a\}, \{a, b\}\}$ and $\{\emptyset, \{b\}, \{a, b\}\}$ would be separate topologies on $\{a, b\}$. Later in the paper, we will refine our representations and only consider the **equivalent**, or **unlabeled**, **topologies** on X . In these instances, the two topologies above would not be considered distinct. As we progress, we will utilize **Hasse diagrams** to explore the unlabeled topologies on a finite set. These directed graphs are a rendering of a partially ordered set whose edge set relation is defined by the cover relation of the partially ordered set. These graphs are **transitive**, meaning that if $x < y$ and $y < z$ for some vertices $x, y, z \in V(G)$ and some cover relation $<$, then $x < z$.

Lastly, we will return to the undirected, loopless graph structure and examine the **degrees of a vertex** v , $d(v)$, or the number of edges incident to v . These graphs will be found to be **k -regular**, meaning every $v \in V(G)$ has degree $d(v) = k$.

3. NUCLEUS-SHELL GRAPHS

Starting my exploration from the graph in Figure 1, I defined my topological graphs by the following rules:

1. For the graph G_X , let $V(G_X)$ be some subset of $P(X) - \{\emptyset, X\}$.
2. Let the edges of $E(G_X)$ be defined by $v_1 \leftrightarrow v_2$ if $v_1 \cup v_2 \cup \{\emptyset, X\}$ is a topology on X , for some $v_1, v_2 \in V(G_X)$.

Note that the question as to which selection of subsets of $P(X)$ which will serve as $V(G_X)$ is a complicated. As shown in the graph in Figure 1, not all of the 256 possible subsets of $P(X)$ are drawn. This issue is discussed in greater detail below.

When $|X| = 1$ or 2, the graphical representations are trivial given our above notational restrictions. They are, however, detailed below. Once $|X| > 2$, a greater organizational structure is needed. For certain vertices, as discussed in the introduction, $v_i \in V(G_X)$, $v_i \cup \{\emptyset, X\}$ is a topology on X . This is not necessarily the case for all vertices however. I defined two concepts in an attempt to simplify this complex problem: the nucleus and shell graphs.

Definition 2. In a G_X , the **nucleus**, $N_{|X|}$, is the full subgraph of G_X in which $V(N_n) = \{A | A \text{ is a non-empty proper subset of } X\}$. Thus $V(N_n)$ consists of all the singleton subsets of $P(X) - \{\emptyset, X\}$. Note that $V(N_n) = V_1(G_n)$, as defined in the introduction.

Example 1. For $n = 3$, $V(N_3) = V_1(G_3) = \{a, b, c, ab, ac, bc\}$. Similarly, for $n = 4$, $V(N_4) = \{a, b, c, d, ab, ac, ad, bc, bd, cd, abc, abd, acd, bcd\}$.

Thus every topology with three or four open sets on a given X_n is represented as a vertex or an edge of N_n . Note that $|V(N_n)| = 2^n - 2$, as these are the number of non-empty, proper subsets of X .

Note that we have an edge $(u, v) \in E(N_{|X|})$ if and only if $u \subset v$, $v \subset u$, or $u = v^c$. Thus we can count the number of edges in N_n . The number of pairs (u, v) such that $u = v^c$ is $\frac{2^n - 2}{2} = 2^{n-1} - 1$. Also, the number of pairs (u, v) with $u \subset v$ is $\sum_{k=2}^{n-1} \binom{n}{k} (2^k - 2) = 3^n - 3 \cdot (2^n) + 3$. Thus the number of edges in N_n is defined by the mapping

$$f(n) = 3^n - 3 \cdot 2^n + 2^{n-1} + 2.$$

Example 2. Let $X_1 = \{a\}$.

Since X_1 's only topologies are the discrete and indiscrete topologies, neither of which are drawn in this method, there is no graphical representation. Since there is no graph, the total number of edges is $f(1) = 0$, to which we add 1 for the discrete and indiscrete topologies, which coincide in this case, giving us the total number of topologies on a set of one element: 1.

Example 3. Let $X_2 = \{a, b\}$.

The graph of topologies on $X_2 = \{a, b\}$ is the isolated vertices a and b . An edge drawn between them would represent the topology $\{\emptyset, \{a\}, \{b\}, \{a, b\}\}$, which is the power set of X_2 and therefore is not pictured. However, each of these vertices would form a topology in and of themselves. This graph is N_2 , the first instance of a nucleus subgraph.

Thus, the topologies on X_2 are formed by each of the two vertices in the nucleus, plus the discrete and indiscrete topologies, for a total count of 4 topologies.

As I began to examine the graph of the topologies, I developed the concepts of the nucleus and shells of the graph in order to better organize my system in anticipation of attempting to graph the topologies of larger sets. I will first show the nucleus, N_3 , then its single shell S_3^2 , and finally the complete graph. It is known that there are 29 labeled topologies on this set [9].

N_3 , pictured in Figure 2, contains $2^n - 2 = 6$ vertices.

We see in the graph of N_3 below the $f(3) = 9$ edges representing the topologies with four sets, six of the edges representing subset relationships between the vertices and three representing vertex pairs which partition X .

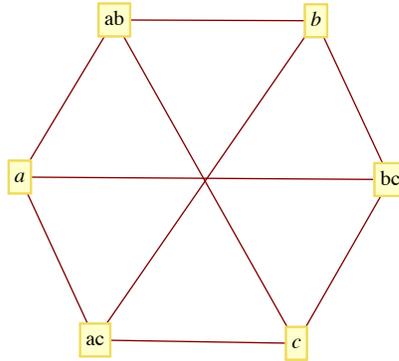


Figure 2: Nucleus Graph, N_3

These values, plus the fact that each vertex in the nucleus is a topology when combined with \emptyset and X_3 , gives a total of 15 distinct topologies.

My plan was to define shells around this nucleus by making the k th shell, S_n^k , consist of some judicious choice of k element subsets of $P(X_n)$. I define S_3^2 to be the pairs of open sets of the same order.

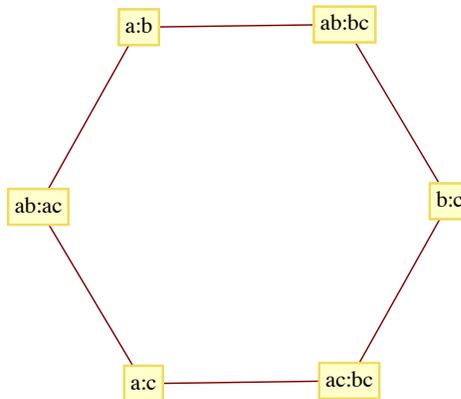


Figure 3: Shell Graph of two open sets, S_3^2

Thus, all topologies of three open sets, such as $\{\emptyset, \{a\}, X_3\}$, are counted by the vertices of the nucleus, all topologies of four open sets are counted by the edges of the nucleus, and all topologies of six open sets are counted by the edges of S_3^2 . Once these two subgraphs are combined together, the edges connecting the nucleus to the shell will give us all the topologies of five open sets.

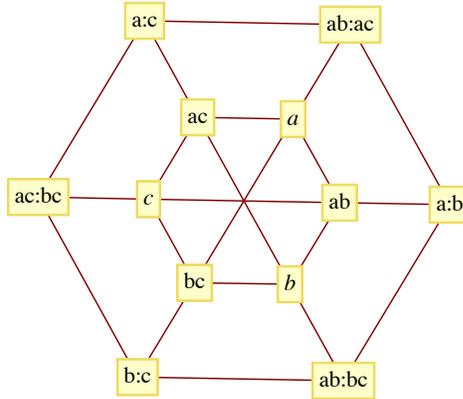


Figure 4: Nucleus-Shell Graph on $X_3 = \{a, b, c\}$

The nucleus contains, as stated above, 15 topologies; S_3^2 contains 6. The two are joined together by 6 five set topologies, such as $a : b \leftrightarrow ab = \{\emptyset, \{a\}, \{b\}, \{a, b\}, X\}$. This gives a total of 27 topologies pictured, plus the discrete and indiscrete topologies, for a total of 29 labeled topologies.

After such a successful representation, I attempted to expand this construction to $X_4 = \{a, b, c, d\}$.

Example 4. In Figure 5, we have the graph N_4 . I chose to render the graph in three dimensions for the sake of clarity. The outer edges in the diagram connect sets u and v with $u \subset v$ and the inner edges connect sets u and v with $u = v^c$. It is also clear that permutations of N_3 is an often repeated subgraph; in fact, every vertex of N_4 lies on at least 2 subgraphs isomorphic to N_3 .

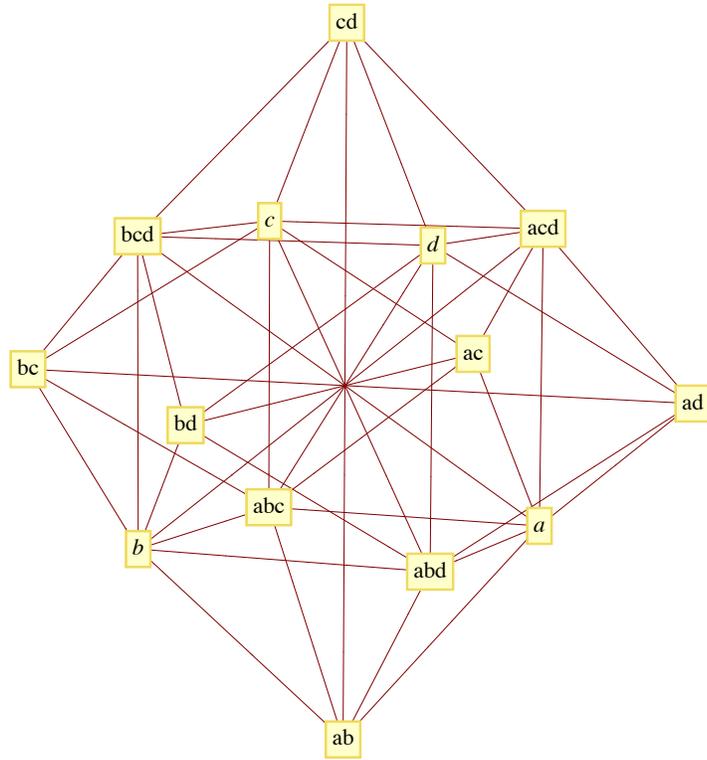


Figure 5: Nucleus Graph, N_4

Thus the nucleus gives a total of 57 labeled topologies on X_4 , still far from the 355 known.

When attempting to construct the first shell, S_4^2 , several difficulties in construction occurred.

Using the same method of choosing subsets of $P(X)$ to serve as the vertices of S_4^2 , I attempted to determine, before construction, the number of vertices in S_4^2 . Since S_4^2 will contain the pairs of elements of the same cardinality, the vertices of S_4^2 are shown in Table 1 below.

Table 1: Vertices of S_4^2

Cardinality of sets	Example	Total
1	$a : b$	6
2	$ab : bc$	12
2	$ab : cd$	3
3	$abc : abd$	6

So, prior to construction of S_4^2 , $V(S_4^2)$ contains 27 vertices. Furthermore, not all topologies on five sets would be captured by this selection of vertices for S_4^2 .

For example, when constructing the topology

$$\{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}, X_4\},$$

which in the Nucleus-Shell system would be an edge connecting a vertex in N_4 to a vertex in S_4^2 , how should the topology be represented? As the edge $a \leftrightarrow ab : abc$, $ab \leftrightarrow a : abc$, or $abc \leftrightarrow a : ab$? It seemed I could not, as I had done before, solely choose vertices of pairs of open sets of the same cardinality.

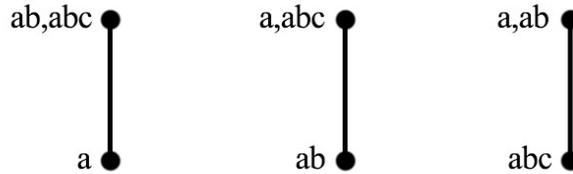


Figure 6: Possible edges of G_{X_4} which form the topology $T = \{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}, X_n\}$

Each choice dramatically altered the overall shape and construction of S_4^2 , as this choice determined the vertices that would be available to construct the topologies of six open sets whose representative edges would lie wholly in S_4^2 . In the pursuit of simplicity, I attempted to minimize the number of vertices needed for S_4^2 . However, in any construction, the shell was not as well organized as the nucleus had been, making the joining of the two graphs,

as had been seen with N_3 and S_3^2 , unenlightening. Another downfall of this construction, which was used initially to count the number of topologies on a three element set, is that for the construction of these graphs for $n \geq 4$ it is necessary to have all the topologies on the set before work begins. No obvious pattern had presented itself to allow me to extrapolate the topologies, or their graph structure, for higher values of n .

With the shell construction abandoned, greater attention was paid to the nucleus graph. With this came the realization that all topologies on X were paths on the nucleus graph and that the shell structure was truly redundant. As seen in Figures 7 and 8, the topology $\{\{a\}, \{b\}, \{a, b\}, \{b, c\}\} \cup \{\emptyset, X_i\}$ corresponds to a path on N_3 and on N_4 , respectively.

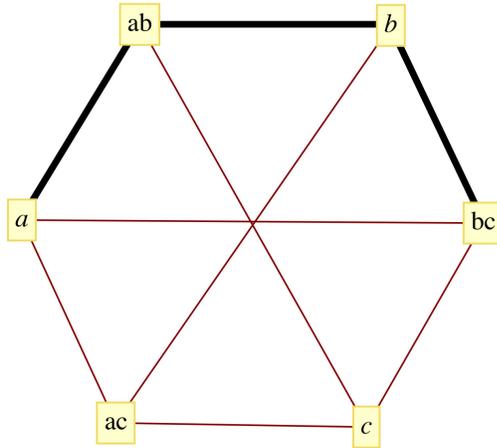


Figure 7: $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}, \{b, c\}, X_3\}$ highlighted on N_3

However, while every topology can be drawn as a path on the nucleus graph, not every path is in turn a topology on X .

Again, it seemed that the nucleus graph, even with the outer shell structure abandoned, would not prove a fruitful path of inquiry. This led to a question: could the topologies on a set of n elements be used to construct the topologies on a set of $n + 1$ elements? I had already seen that each topology was the union of smaller topologies on the set, or the union of edges on a graph forming a path, and that any topology on n elements had a

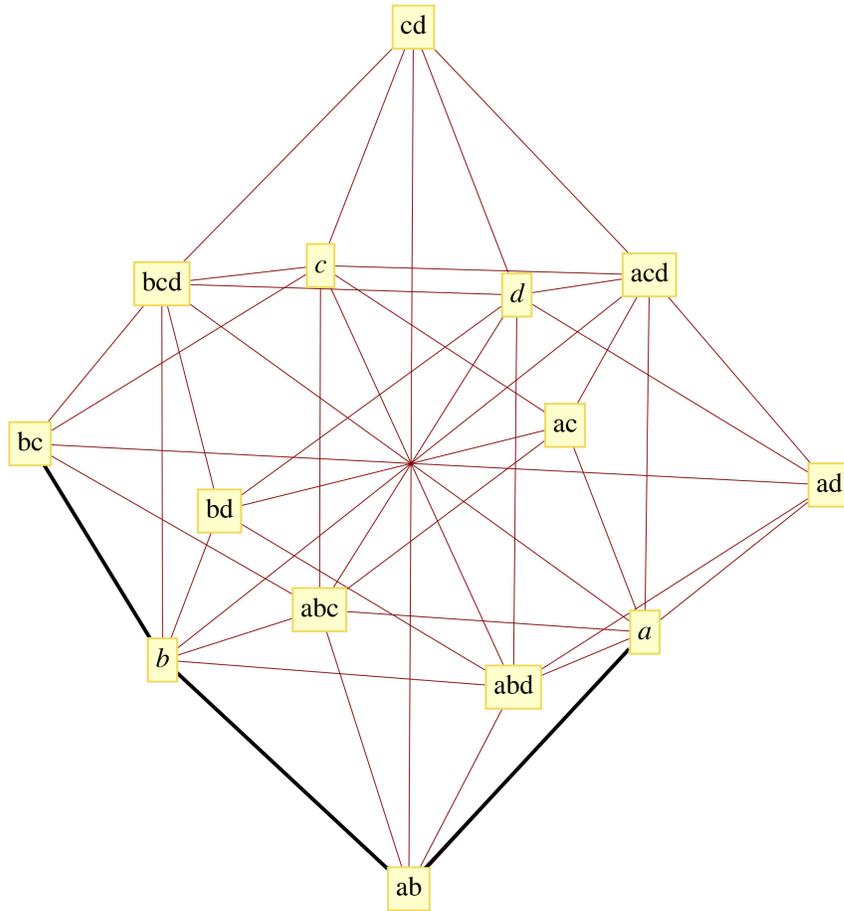


Figure 8: $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}, \{b, c\}, X_4\}$ highlighted on N_4

correlate on $n + 1$ elements, as seen in Figures 7 and 8. This led to my first attempt at a recursive definition for the construction of the topologies on $n + 1$ elements from the topologies on an n element set.

4. A RECURSIVELY DEFINED SET OF TOPOLOGIES

In the nucleus and shell graphs, needing to have all the topologies on the finite set proved to be a major hurdle in representing the topologies graphically. Finding all the topologies on a set through a brute force method is already tedious on a set of 4 elements, and quickly became infeasible as the number of elements increases. For these reasons, my research focus shifted to trying to find a means of generating the topologies on a set of $n + 1$ elements from the topologies on a set of n elements. Working with the nucleus-shell graphs did present some interesting relationships between topologies on sets differing by one element.

Let X_n be a set with n elements, let $x \notin X_n$ and let $X_{n+1} = X_n \cup \{x\}$. Then for any topology $\{\emptyset, A_1, \dots, A_k, X_n\}$, where each A_i is a non-empty, proper subset of X_n , the set $\{\emptyset, A_1, \dots, A_k, X_n, X_{n+1}\}$ is a topology on a set of $n + 1$ elements. Also, $\{\emptyset, \{x\}, A_1 \cup \{x\}, \dots, A_k \cup \{x\}, X_{n+1}\}$ is a topology on X_{n+1} . If for all sets A_i, A_j in a topology on X_n we have that $A_i \cap A_j \neq \emptyset$, then $\{\emptyset, A_1 \cup \{x\}, \dots, A_k \cup \{x\}, X_{n+1}\}$ is a topology on the set X_{n+1} as well. Given these realizations, the following definitions are made.

Definition 3. Let X_n be a finite set, T a topology on X_n , and A_i is a non-empty proper subset of X_n in T . Define the sets xT , $\hat{x}T$, and $x'T$, for some singleton $x \notin X_n$ as follows:

$$xT = \{A_i \cup \{x\} \mid A_i \in T\} \cup \{\emptyset, X_{n+1}\},$$

$$x'T = xT \setminus \{x\},$$

$$\hat{x}T = T \cup \{X_{n+1}\}.$$

It is necessary to make further distinctions on the type of sets which compose the topology

T . Whether or not these classify topologies is dependent on the presence of non-empty open sets with an empty intersection.

Definition 4. T is a **conjoint** topology on X means that for all non-trivial proper subsets A and B of X , if A and B are in T then $A \cap B \neq \emptyset$.

Proposition 1. *If T is a topology on X_n , then we have the following:*

i) xT is a topology on X_{n+1} ,

ii) $\hat{x}T$ is a topology on X_{n+1} ,

iii) $x'T$ is a topology on X_{n+1} if and only if T is conjoint,

iv) $xT \cup \hat{x}T$ is a topology on X_{n+1} ,

v) $x'T \cup \hat{x}T$ is a topology on X_{n+1} if and only if T is conjoint.

Proof. Let X_n be a finite set, $x \notin X_n$, such that $X_{n+1} = X \cup \{x\}$.

Note first that $\{\emptyset, X_{n+1}\}$ is a subset of $xT, x'T$, and $\hat{x}T$, by construction, and as such, $\{\emptyset, X_{n+1}\}$ is a subset of any possible unions of these sets. Thus, we need only check for each case closure by set union and intersection for each set.

i) Let $B, B' \in xT$. Then

$$\begin{aligned} B \cup B' &= (A_i \cup \{x\}) \cup (A_j \cup \{x\}) \text{ for some } A_i, A_j \in T \\ &= (A_i \cup A_j) \cup \{x\} \end{aligned}$$

and

$$\begin{aligned} B \cap B' &= (A_i \cup \{x\}) \cap (A_j \cup \{x\}) \text{ for some } A_i, A_j \in T \\ &= (A_i \cap A_j) \cup \{x\}. \end{aligned}$$

Since $A_i \cup A_j, A_i \cap A_j \in T$, then $(A_i \cup A_j) \cup \{x\}$ and $(A_i \cap A_j) \cup \{x\} \in xT$. Thus xT is closed under unions and intersections. Thus xT is a topology on X_{n+1} .

ii) Now from the construction of $\hat{x}T$, it is clear that for any topology T on X_n , $T \cup \{X_{n+1}\}$ is a topology on X_{n+1} .

iii) By a similar argument as had been made for xT , the set $x'T$ is closed under unions and intersections. It need only be shown that in a conjoint topology, there are no sets in $x'T$ such that their intersection is $\{x\}$.

Assume that there are sets C, C' in $x'T$ such that $C \cap C' = \{x\}$. Then

$C = A_i \cup \{x\}, C' = A_j \cup \{x\}$, where $A_i, A_j \neq \emptyset$. Then we have that

$(A_i \cup \{x\}) \cap (A_j \cup \{x\}) = (A_i \cap A_j) \cup \{x\}$. So $C \cap C' = (A_i \cap A_j) \cup \{x\} = \{x\}$. However, $A_i \cap A_j \neq \emptyset$, since T is a conjoint topology. Thus $C \cap C' \neq \{x\}$. So $x'T$ is a topology on $X_n \cup \{x\} = X_{n+1}$.

iv) Let $D, D' \in xT \cup \hat{x}T$.

If $D, D' \in xT$ or $D, D' \in \hat{x}T$, then, as shown above, the set $xT \cup \hat{x}T$ is closed under unions and intersections. Assume, without loss of generality, that $D \in xT, D' \in \hat{x}T$. Then

$D = A_i \cup \{x\}$ and $D' = A_j$, for some $A_i, A_j \in T$. So we have that

$$\begin{aligned} D \cup D' &= (A_i \cup \{x\}) \cup A_j \\ &= (A_i \cup A_j) \cup \{x\} \\ &\in xT \\ &\subset xT \cup \hat{x}T \end{aligned}$$

and

$$\begin{aligned} D \cap D' &= (A_i \cup \{x\}) \cap A_j \\ &= (A_i \cap A_j) \cup (\{x\} \cap A_j) \\ &= A_i \cap A_j \\ &\in \hat{x}T \\ &\subset xT \cup \hat{x}T \end{aligned}$$

so $xT \cup \hat{x}T$ is a topology on $X_n \cup \{x\} = X_{n+1}$.

v) By similar argument, we have $x'T \cup \hat{x}T$ is a topology on X_{n+1} as long as T is conjoint, shown in case (ii).

□

We now investigate whether every topology on an $n + 1$ element set can be obtained by applying of the the operations $T \rightarrow xT, T \rightarrow x'T, T \rightarrow \hat{x}T, T \rightarrow xT \cup \hat{x}T$, or $T \rightarrow x'T \cup \hat{x}T$ to an n element set. Letting \mathcal{T}_{X_n} be the collection of all topologies on X_n and $T_{n,k}$ be the k th topology on X_n , the following examples will serve to better illuminate the sets defined above.

Example 5. Let $X_0 = \emptyset$. Then the only topology on X_0 is $T_{0,1} = \{\emptyset\}$. Let $X_1 = \{a\}$. Then the only topology on X_1 is $T_{1,1} = \{\emptyset, \{a\}\}$. We note that $T_{1,1} = aT_{0,1}$.

Then $\mathcal{T}_{X_1} = \{T_{1,1}\}$.

Choose $b \notin X$. Then we have

$$\begin{aligned}
 bT_{1,1} &= \{\emptyset, \{b\}, \{a,b\}\} \\
 \hat{b}T_{1,1} &= \{\emptyset, \{a,b\}\} \\
 b'T_{1,1} &= \{\emptyset, \{a\}, \{a,b\}\} \\
 bT_{1,1} \cup \hat{b}T_{1,1} &= \{\emptyset, \{b\}, \{a,b\}\} \\
 bT_{1,1} \cup b'T_{1,1} &= \{\emptyset, \{a\}, \{b\}, \{a,b\}\}.
 \end{aligned}$$

Let $X_2 = \{a, b\}$. The topologies on X_2 , and their expressions using the above operations, are as follows:

$$\begin{aligned}
 T_{2,1} &= \{\emptyset, \{a, b\}\} = \hat{b}T_{1,1} \\
 T_{2,2} &= \{\emptyset, \{a\}, \{a, b\}\} = b'T_{1,1} \\
 T_{2,3} &= \{\emptyset, \{b\}, \{a, b\}\} = bT_{1,1} \cup \hat{b}T_{1,1} \\
 T_{2,4} &= \{\emptyset, \{a\}, \{b\}, \{a, b\}\} = bT_{1,1} \cup b'T_{1,1}.
 \end{aligned}$$

Note that $bT_{1,1} \cong b'T_{1,1} \cong bT_{1,1} \cup \hat{b}T_{1,1}$.

Applying our operation to the labeled topologies on $\{a, b\}$ we obtain the following 17 labeled topologies on $\{a, b, c\}$.

Table 2: Topologies generated from X_2

	$T_{2,1} = \{\emptyset, \{a, b\}\}$	$T_{2,2} = \{\emptyset, \{a\}, \{a, b\}\}$	$T_{2,4} = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$
$cT_{i,j}$	$\{\emptyset, \{c\}, \{a, b, c\}\}$	$\{\emptyset, \{c\}, \{a, c\}, \{a, b, c\}\}$	$\{\emptyset, \{c\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$
$c'T_{i,j}$	$\{\emptyset, \{a, b, c\}\}$	$\{\emptyset, \{a, c\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{b\}, \{a, b\}, \{a, b, c\}\}$
$\hat{c}T_{i,j}$	$\{\emptyset, \{a, b\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}\}$	$\{\emptyset, \{a, c\}, \{b, c\}, \{a, b, c\}\}^*$
$cT_{i,j} \cup \hat{c}T_{i,j}$	$\{\emptyset, \{c\}, \{a, b\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{c\}, \{a, b\}, \{a, c\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$
$c'T_{i,j} \cup \hat{c}T_{i,j}$	$\{\emptyset, \{a, c\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{a, b\}, \{a, c\}, \{a, b, c\}\}$	$\{\emptyset, \{a\}, \{b\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}^*$

* Note that $\hat{c}T_{2,4}$ and $c'T_{2,4} \cup \hat{c}T_{2,4}$ are not topologies on $X_3 = \{a, b, c\}$ since $T_{2,3}$ is a disjoint topology.

1. $\{\{\emptyset\}, \{a, b, c\}\}$
2. $\{\{\emptyset\}, \{c\}, \{a, b, c\}\}$
3. $\{\{\emptyset\}, \{a, b\}, \{a, b, c\}\}$
4. $\{\{\emptyset\}, \{a, c\}, \{a, b, c\}\}$
5. $\{\{\emptyset\}, \{b, c\}, \{a, b, c\}\}$
6. $\{\{\emptyset\}, \{a\}, \{a, b\}, \{a, b, c\}\}$
7. $\{\{\emptyset\}, \{b\}, \{a, b\}, \{a, b, c\}\}$
8. $\{\{\emptyset\}, \{c\}, \{a, b\}, \{a, b, c\}\}$
9. $\{\{\emptyset\}, \{c\}, \{a, c\}, \{a, b, c\}\}$
10. $\{\{\emptyset\}, \{c\}, \{b, c\}, \{a, b, c\}\}$
11. $\{\{\emptyset\}, \{a\}, \{b\}, \{a, b\}, \{a, b, c\}\}$
12. $\{\{\emptyset\}, \{a\}, \{a, b\}, \{a, c\}, \{a, b, c\}\}$
13. $\{\{\emptyset\}, \{b\}, \{a, b\}, \{b, c\}, \{a, b, c\}\}$
14. $\{\{\emptyset\}, \{c\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$
15. $\{\{\emptyset\}, \{a\}, \{c\}, \{a, b\}, \{a, c\}, \{a, b, c\}\}$
16. $\{\{\emptyset\}, \{b\}, \{c\}, \{a, b\}, \{b, c\}, \{a, b, c\}\}$
17. $\{\{\emptyset\}, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$

This list contains at least one copy of each of the 9 inequivalent topologies possible on a 3

element set, which are listed below.

$$\begin{aligned}
T_{3,1} &= \{\emptyset, \{a, b, c\}\} = c'T_{2,1} \\
T_{3,2} &= \{\emptyset, \{c\}, \{a, b, c\}\} = cT_{2,1} \\
T_{3,3} &= \{\emptyset, \{a, c\}, \{a, b, c\}\} = c'T_{2,2} \\
T_{3,4} &= \{\emptyset, \{c\}, \{a, c\}, \{a, b, c\}\} = cT_{2,2} \\
T_{3,5} &= \{\emptyset, \{a\}, \{b\}, \{a, b\}, \{a, b, c\}\} = c'T_{2,4} \\
T_{3,6} &= \{\emptyset, \{a\}, \{a, b\}, \{a, c\}, \{a, b, c\}\} = c'T_{2,2} \cup \hat{c}T_{2,2} \\
T_{3,7} &= \{\emptyset, \{a\}, \{c\}, \{a, b\}, \{a, c\}, \{a, b, c\}\} = cT_{2,2} \cup \hat{c}T_{2,2} \\
T_{3,8} &= \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\} = cT_{2,4} \cup \hat{c}T_{2,4} \\
T_{3,9} &= \{\emptyset, \{c\}, \{a, b\}, \{a, b, c\}\} = cT_{2,1} \cup \hat{c}T_{2,1}
\end{aligned}$$

To obtain all the equivalent topologies on $\{a, b, c\}$, we could apply our operations to $\{a, c\}$ using b , and to $\{b, c\}$ using a .

From this work I hoped to prove that all topologies on a set X_{n+1} were in fact one of the above defined sets derived from a topology on X_n . However, a counter-example was found when examining the recursion from X_3 to X_4 .

Let $X_4 = \{a, b, c, d\}$, $X_3 = \{a, b, c\}$ and consider the topology $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}, X_4\}$ on X_4 . Let T' be some topology on X_3 .

Since $\{a\} \in T$ we cannot have $T = dT'$ or $T = d'T'$. Since $X_3 \notin T$, and $X_3 \in \hat{d}T'$ for all T' , we cannot have $T = \hat{d}T'$, $T = dT' \cup \hat{d}T'$ or $T = d'T' \cup \hat{d}T'$.

This construction would not ultimately work to produce all of the topologies on X_{n+1} from the topologies on the set X_n .

However, this counter-example provided insight into the way that a topology can be broken down into topologies on a set of cardinality one less. Rather than working from X_n

to X_{n+1} , the opposite direction was pursued in order to determine if there was some pattern to why the recursive definition had failed.

5. PARTITIONING TOPOLOGIES

As the recursive definition of the last section failed to capture all topologies, we study further how topologies on an $n + 1$ element set are related to topologies on an n element subset. The goal is to obtain insight into how to make a recursive definition for producing all the finite topologies.

Given any topology T on a set X and an element $a \in X$ there are two natural ways to obtain a topology on the set $X - \{a\}$. Observe that the sets in T which contain a are closed under unions and intersections, so those sets together with \emptyset form a sub-topology on X and we can obtain a topology on $X - \{a\}$ by deleting a from every set in the sub-topology. Similarly we note that the sets in T which do not contain a are also closed under unions and intersections, so by taking these sets together with $X - \{a\}$ we obtain a topology on $X - \{a\}$.

More formally, we define three sets $S_a, S_{\bar{a}}$, and \bar{S}_a as follows: $S_a = \{A_i \in \bar{T} \mid a \in A_i\}$, $S_{\bar{a}} = \{A_i \in \bar{T} \mid a \notin A_i\}$ and finally $\bar{S}_a = \{A_i - \{a\} \mid A_i \in S_a\}$.

Definition 5. Let $\mathbf{T}_a = \bar{S}_a \cup \{\emptyset, X\}$ and $\mathbf{T}_{\bar{a}} = S_{\bar{a}} \cup \{\emptyset, X\}$.

We record our initial observations in the following proposition.

Proposition 2. *Let T be a topology on a set X and let $a \in X$. Then T_a and $T_{\bar{a}}$ are topologies on $X - \{a\}$.*

Example 6. Let $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}, \{a, c, d\}, X_4\}$. For each choice of an element of X_4 , the partitioning can produce different (though not necessarily inequivalent) topologies on X_3 , some three element set, as shown below:

$$T_a = \{\emptyset, \{b\}, \{cd\}, X_3\} \text{ for } X_3 = \{b, c, d\}$$

$$T_{\bar{a}} = \{\emptyset, \{b\}, X_3\}$$

$$T_b = \{\emptyset, \{a\}, X_3\} \text{ for } X_3 = \{a, c, d\}$$

$$T_{\bar{b}} = \{\emptyset, \{a\}, X_3\}$$

$$T_c = \{\emptyset, \{a, d\}, X_3\} \text{ for } X_3 = \{a, b, d\}$$

$$T_{\bar{c}} = \{\emptyset, \{a\}, \{b\}, \{a, b\}, X_3\}$$

$$T_d = \{\emptyset, \{a, c\}, X_3\} \text{ for } X_3 = \{a, b, c\}$$

$$T_{\bar{d}} = \{\emptyset, \{a\}, \{b\}, \{a, b\}, X_3\}.$$

It is easy to see that $T_{\bar{a}} \cong T_b \cong T_{\bar{b}}$, $T_c \cong T_d$, and $T_{\bar{c}} \cong T_{\bar{d}}$.

Next we examined whether there were certain properties of T which would be preserved by the partitioning process. Most notably the T_0 property was preserved.

Lemma 1. *Every T_0 topology contains a singleton set.*

Proof. Let T be a T_0 topology on a set X and assume that T does not have a singleton set.

Let $A, |A| \geq 2$, be the non-empty set in T with the fewest number of elements.

Let $a_1, a_2 \in A, a_1 \neq a_2$.

Then there exists a $B \in T$ such that, without loss of generality, $a_1 \in B$ and $a_2 \notin B$ since T is a T_0 topology.

Then $a_1 \in A \cap B$ and $a_2 \notin A \cap B$. This implies that $|A \cap B| < |A|$, which is a contradiction.

Thus, T contains a singleton set.

□

Proposition 3. *If T is a T_0 topology on X_{n+1} , then for some $a \in X_{n+1}$, T_a is a T_0 topology on X_n .*

Proof. Let T be a T_0 topology on X_{n+1} .

By Lemma 1, there exists an $a \in X_{n+1}$ such that $\{a\} \in T$. Then we have from Proposition 2 we have that T_a is a topology on X_n . It will suffice to show that T_a is T_0 on X_n .

Let $u, v \in X_n$. Then u, v are distinct points with respect to T . So, there exists an open set $U \in T$ such that, without loss of generality, $u \in U, v \notin U$. Since $\{a\} \in T$ then $W = U \cup \{a\} \in T_a$ and $u \in W - \{a\}$ and $v \notin W - \{a\}$.

Thus T_a is a T_0 topology on $X_{n+1} - \{a\} = X_n$.

□

By a simple counter-example, however, we see that $T_{\bar{a}}$ is not necessarily T_0 if T is. Letting $T = \{\emptyset, \{a\}, \{a, b\}, \{b, c\}, \{a, b, c\}, \{b, c, d\}, X_4\}$ we have that $T_{\bar{a}} = \{\emptyset, \{b, c\}, X_3\}$, which is not a T_0 topology on $X_3 = \{b, c, d\}$. However, in the case of the topology $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}, \{b, c\}, X_3\}$, $T_{\bar{a}} = \{\emptyset, \{b\}, \{b, c\}, X_3\}$ is a T_0 topology. Ultimately, the conditions under which $T_{\bar{a}}$ preserved the T_0 property led to the definition of a maximal, complete chain, which is explored in the following section and became the central focus of my research from this point forward.

I turned to contemporary research to try and help elucidate some of the reasons why the construction worked for some topologies and not for others. In [5], Ern  showed that $T(n)$, the number of topologies on a set of n elements, is asymptotically equal to $T_0(n)$, the number of T_0 topologies on n elements. For smaller n , Evans et al. [6] and Renteln [10] reduced the computation of $T(n)$ to the number of partially ordered set on n elements. These values are connected by the following formula:

Let $T(n)$ denote the number of distinct topologies on a set with n points. The number of distinct T_0 topologies on a set with n points, denoted $T_0(n)$, is related to $T(n)$ by the formula

$$T(n) = \sum_{k=0}^n S(n,k)T_0(k)$$

where $S(n,k)$ denotes the Stirling number of the second kind.

Since all the topologies on X_{n+1} are enumerated in terms of all the T_0 topologies on $X_0 \dots X_n$ I decided to focus on just the T_0 topologies.

6. MAXIMAL COMPLETE CHAINS

Solely considering T_0 topologies on a finite set, I examine their Hasse diagrams.

Uniformly, these Hasse diagrams contained a nested collection of sets whose order increased by one at each level of the diagram, from \emptyset to X_n . It was this feature, seen in Figure 10 and 11, which is defined as a maximal, complete chain (or MC-chain).

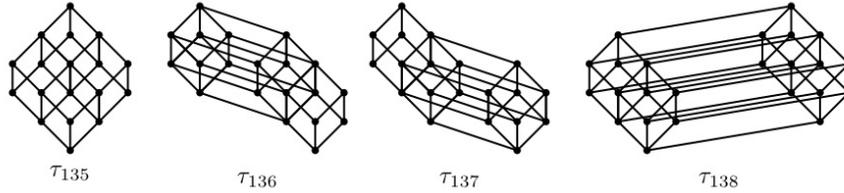


Figure 9: Sample Hasse diagrams of topologies on the set X_5 , [4]

Definition 6. Let $T = \{\emptyset, A_1, A_2, \dots, A_{n-1}, X_n\}$ be a topology on X_n . Then T is a **maximal, complete chain** (or **MC-chain**) if for all $A_i \in T$, $|A_i| = i$, and $A_i \subsetneq A_{i+1}$, where $A_0 = \emptyset$ and $A_n = X_n$.

For example, on the set $X = \{a, b, c, d\}$, the set $\{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}, \{a, b, c, d\}\}$ is a maximal, complete chain.

Example 7. Pictured below are two more examples of MC-chains, highlighted in the Hasse diagrams of two T_0 topologies on X_5 and X_6 . The first represents the unlabeled T_0 topology

$$\{\emptyset, \{a\}, \{a, b\}, \{a, c\}, \{a, d\}, \{a, b, c\}, \{a, b, d\}, \{a, c, d\}, \{a, b, c, d\}, \{a, b, c, d, e\}\}$$

with the MC-chain $\{\emptyset, \{a\}, \{a, b\}, \{a, b, c\}, \{a, b, c, d\}, \{a, b, c, d, e\}\}$ highlighted. In each Figure, one can see that there are in fact several possible MC-chains in each T_0 topology. This property of MC-chains will be discussed more fully later in the section.



Figure 10: Hasse diagram of a T_0 topology on X_5 , MC-chain highlighted



Figure 11: Hasse diagram of a T_0 topology on X_6 , MC-chain highlighted

Lemma 2. *All maximal complete chains are T_0 topologies.*

Proof. Let T be an MC-chain on the set X_n . Then $\emptyset, X_n \in T$. Also, for any $A_i, A_j \in T$, $A_i \subsetneq A_j$ or $A_j \subsetneq A_i$. Without loss of generality, assume the former.

Then $A_i \cup A_j = A_j$ and $A_i \cap A_j = A_i$, both sets in T . Thus, T is closed under unions and intersections.

Finally, assume that for some $x, y \in X_n$, $x, y \notin A_i$ and $x, y \in A_{i+1}$. However, this would imply that $|A_i| + 2 = |A_j|$, in which case T is not an MC-chain. Thus, for any $x, y \in X_n$, there exists some A_{i+1} such that $x \in A_{i+1}, y \notin A_{i+1}$.

Thus, T is a T_0 topology.

□

Remark 1. An MC-chain is a minimal T_0 topology on a set of n elements.

Letting T be a proper subset of some MC-chain \mathcal{A} , we note that if we omit any set $A_i \in \mathcal{A}$, then the two elements in $A_{i+1} - A_{i-1}$ are topologically indistinguishable.

It should be noted that similar constructions have been used in modern research in finite set topology. In [1], the Adamenko and Velichko use a topological quiver, or T-quiver, from \emptyset to X_n . While slightly different in its construction from an MC-chain, the T-quiver serves the same purpose - creating a path of open sets in a T_0 space. The authors asserts that “[a]t the k th level of an arbitrary T-quiver of T_0 -topology, there are vertices corresponding to k -element open sets”, which is, in terms of MC-chains, the fact that $|A_i| = i$. The major difference between the T-quivers and MC-chains is only that some union of MC-chains would make up a T-quiver. They go on to prove that for any topology on an n -element set, it can only be T_0 if its corresponding T-quiver has n levels. Adamenko considers the empty set as the zero level of the T-quiver, so this statement is equivalent to Lemma 3 below, that T is T_0 if it contains an MC-chain.

Lemma 3. *If T is T_0 on X_n , then T contains a maximal complete chain.*

Proof. Proof by induction.

As $n = 1$ is a trivial case, we start with $n = 2$.

So $X_2 = \{x_1, x_2\}$. There are two inequivalent T_0 topologies on X :

$$\{\emptyset, \{x_1\}, \{x_1, x_2\}\}$$

$$\{\emptyset, \{x_1\}, \{x_2\}, \{x_1, x_2\}\}$$

The first is a maximal, complete chain (MC-chain) and is also a subset of the second. So for our base case the proposition holds.

Assume that for all $k < n$ the proposition holds where $n > 2$.

Let T be a T_0 topology on X_n .

By Lemma 1, T contains a singleton set $\{x_1\}$.

Let $Y = X_n - \{x_1\}$ and define $T_x = \{A - \{x_1\} | A \in T\}$.

Since T is T_0 on X , T_x is T_0 on $X - \{x_1\} = Y$.

Since T_x is T_0 on a set Y of $n - 1$ elements and $n - 1 \geq 2$, T_x contains an MC-chain $\emptyset : x_2 : x_2x_3 : \dots : x_2\dots x_n$ on Y , by the induction hypothesis.

This give us the MC-chain $\emptyset : x_1 : x_1x_2 : x_1x_2x_3 : \dots : x_1x_2\dots x_n$ as a subset of T on X_n .

□

So we have that each T_0 topology contains at least one MC-chain. Examination of the Hasse diagrams shows that in fact all sets of a T_0 topology are a member set of an MC-chain.

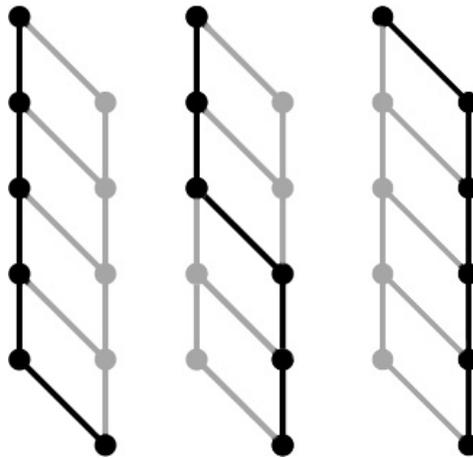


Figure 12: T_0 topology on X_5 , some distinct MC-chains highlighted

Theorem 1. *Every open set in a T_0 topology is an element of an MC-chain.*

Proof. Let T be a T_0 topology on X_n .

Then there exists an MC-chain $\mathcal{A} = \{\emptyset, A_1, A_2, \dots, A_{n-1}, X\}$ that is a subset of T .

Let $B \in T$ such that $B \notin \mathcal{A}$.

Then there exists an $A_i \in \mathcal{A}$ such that $B \subset A_i$ and an $A_j \in \mathcal{A}$ such that $A_j \subset B$, even if these are trivially \emptyset and X , respectively. Thus, we have already the following subsets of \mathcal{A} as parts of our MC-chain which contains B : $\{\emptyset, \dots, A_j\}$ and $\{A_i, \dots, X\}$.

Note that $B \cap A_i = B$.

Consider the sequence of sets

$\emptyset, A_1 \cap B, A_2 \cap B, \dots, A_{n-1} \cap B, B, B \cup A_1, B \cup A_2, \dots, B \cup A_{n-1}, X_n$. Each set in the sequence is either equal to the previous one, or is the union of the previous set with a singleton. Therefore if we remove the duplicates, we obtain an MC-chain containing B .

□

Since every T_0 topology T contains an MC-chain, these are the only minimal T_0 topologies. Also, we have that every set in T is in some MC-chain, which is visually apparent in the Hasse diagrams as each node lies on a direct path in the diagram of length $n + 1$ from the empty set to X . Thus we may say the following.

Remark 2. Every T_0 topology is a union of MC chains.

It should be noted that the converse doesn't hold. Consider the MC-chains on X_3 :

$\mathcal{A} = \{\emptyset, \{a\}, \{a, b\}, X_3\}$ and $\mathcal{B} = \{\emptyset, \{c\}, \{b, c\}, X_3\}$. Their union is not even a topology on X_3 , as it lacks the open sets $\{\{b\}, \{a, c\}\}$.

Once we have that each T_0 topology is the union of MC-chains, and since the set $P(X)$, the power set of X , is T_0 on X , $P(X)$ is the union of, specifically, $n!$ distinct MC-chains. By “distinct”, we mean that for two MC-chains \mathcal{A}, \mathcal{B} , there exists at least one set $B \in \mathcal{B}$ such that $B \notin \mathcal{A}$. However, the intersection of the MC-chains need not be empty.

The question then became whether or not there was some method in which these MC-chains could be used to form the topologies on a set X_n . The development of MC-chains was for the purpose of finding a recursive definition with which the topologies on a set could be produced. If each set in a T_0 topology is part of an MC-chain, could a relationship between these MC-chains be found for the purposes of recursion? It became necessary to make a more rigorous definition for the specific relationship between the MC-chains which whose union forms a T_0 topology.

Definition 7. Two MC-chains \mathcal{A}, \mathcal{B} are **adjacent** if they differ by only one set.

Note that given any topology T , Definition 7 defines a graph structure G on the set of all MC-chains of possible on a set X . When $|X| = n$ the graph has $n!$ vertices because MC-chains are in a one-to-one correspondence with the ways to arrange the elements of X in order. Also, two MC-chains are adjacent if and only if the corresponding orderings of X are related by swapping two adjacent spots in the ordering. So the graph has $n!$ vertices and is $(n - 1)$ -regular. Given a topology T on X we consider the full sub-graph G_T consisting of the MC-chains and adjacency edges which are in the topology T .

Theorem 2. *For any topology T on a finite set X , the graph G_T is connected.*

Proof. Let $\mathcal{A} \subsetneq T$ be an MC-chain, where T is a non-minimal T_0 topology on X_n . Since T is non-minimal, there exists a set $B \in T$ such that $B \notin \mathcal{A}$. By Theorem 1, let $B \in \mathcal{B}$, an MC-chain in T . Thus \mathcal{A}, \mathcal{B} are distinct vertices of G_T .

Let $A_i \in \mathcal{A}$ and $B_i \in \mathcal{B}$, where i is the smallest value such that $A_i \neq B_i$. We note that $|A_i| = |B_i| = i$.

We will consider this problem in cases.

First, assume $A_i \cup B_i = A_{i+1}$. Then $\mathcal{C} = \{\emptyset, A_1, A_2, \dots, A_{i-1}, B_i, A_{i+1}, \dots, A_{n-1}, X_n\}$ is an adjacent MC-chain to \mathcal{A} on the graph G_n .

Now, assume $A_i \cup B_i \neq A_{i+1}$. Note that there exists some smallest possible set $A_{i+j} \in \mathcal{A}$ such that $B_i \subset A_{i+j}$. So $\mathcal{B} - \mathcal{A} = \{B_i, B_{i+1}, \dots, B_{i+j-1}\}$, and thus has order j . Then consider the MC-chain $\mathcal{C}_1 = \{\emptyset, A_1, \dots, A_i, A_i \cup B_i, A_{i+1} \cup B_i, \dots, A_{i+j-1} \cup B_i, A_{i+j}, \dots, X_n\}$. The set $A_{i+j-1} \cup B_i$ is the last set of \mathcal{C}_1 which is not also an element of \mathcal{A} . So $\mathcal{C}_1 - \mathcal{A} = \{A_i \cup B_i, A_{i+1} \cup B_i, \dots, A_{i+j-1} \cup B_i\}$, and $|\mathcal{C}_1 - \mathcal{A}| = j - 1$. Thus the process has produced an MC-chain \mathcal{C}_1 which has one few set of difference from \mathcal{A} than \mathcal{B} has.

Continue this construction recursively, taking the first set of \mathcal{C}_1 which is not in \mathcal{A} and taking its union with all the sets of higher order in $\mathcal{A} - \mathcal{C}_1$ to produce \mathcal{C}_2 . Thus we have the following sequence of MC-chains:

$$\begin{aligned}
\mathcal{B} &= \{\emptyset, A_1, \dots, B_i, \dots, B_{i+j-1}, A_{i+j}, \dots, A_{n-1}, X_n\} \\
\mathcal{C}_1 &= \{\emptyset, A_1, \dots, A_i, A_i \cup B_i, \dots, A_{i+j-1} \cup B_i, A_{i+j}, \dots, A_{n-1}, X_n\} \\
\mathcal{C}_2 &= \{\emptyset, A_1, \dots, A_i, A_{i+1}, A_{i+1} \cup B_i, \dots, A_{i+j-1} \cup B_i, A_{i+j}, \dots, A_{n-1}, X_n\} \\
&\vdots \\
\mathcal{C}_k &= \{\emptyset, A_1, \dots, A_{i+j-1}, A_{i+j-1} \cup B_i, A_{i+j}, \dots, A_{n-1}, X_n\} \\
\mathcal{A} &= \{\emptyset, A_1, \dots, A_{i+j-1}, A_{i+j}, A_{i+j+1}, \dots, A_{n-1}, X_n\}
\end{aligned}$$

for some $k \in \mathbb{N}$. This sequence corresponds to a path on G_{X_n} from \mathcal{A} to \mathcal{B} .

□

Thus we have arrived at the fact that all T_0 topologies correspond to a connected subgraph of G_n . With this in mind, we revisit the notion of a recursive definition. However, our focus remains on T_0 topologies.

7. RECURSIVE DEFINITION OF T_0 TOPOLOGIES

In this section we show how an extension of the operations considered in Chapter 4 will lead to generating all of the T_0 topologies on a set X with $n + 1$ elements from those on a set with n elements.

We know every topology on an $n + 1$ element set contains the MC-chain

$\emptyset : x_1 : \dots : x_1 \dots x_n : x_1 \dots x_n x_{n+1}$. Given some T_0 topology T on X_{n+1} , let

$T^* = \{A \in T \mid x_{n+1} \notin A\}$. Then T^* is a topology on X_n .

We note that in fact T^* is T_0 on X_n . If $u \neq v$ are in X_n , then there exists sets $A, B \in T$ such that $u \in A, v \notin A$ or $u \notin A, v \in A$. Without loss of generality, assume the former. But then $A_1 = A \cap x_1 x_2 \dots x_n \in T^*$ and $u \in A_1$ and $v \notin A_1$.

Now suppose that U and V are in T^* and $U \cup \{x_{n+1}\}, V \cup \{x_{n+1}\} \in T$. Then

$(U \cap V) \cup \{x_{n+1}\} = (U \cup \{x_{n+1}\}) \cap (V \cup \{x_{n+1}\})$ is in T also. So the collection of sets

$U \in T^*$ such that $U \cup \{x_{n+1}\} \in T$ is closed under intersection. We obtain a minimal such set $W \in T^*$ by taking the intersection of all the sets with this property.

We also note that if we start with any T_0 topology T^* on X_n , such that $x_{n+1} \notin X_n$ and $W \in T^*$, then $T = T^* \cup \{A \cup \{x_{n+1}\} \mid A \in T^* \text{ and } W \subset A\}$ is a T_0 topology on X_{n+1} .

In terms of the Hasse diagrams, we may interpret this as follows. We take the Hasse diagrams of the inequivalent topologies on a set with n elements. In each diagram select a vertex v . We take the set U of all vertices u such that $v \leq u$. For each $u \in U$, we introduce a new vertex u' , which represents the addition of $\{x_{n+1}\}$ to the set represented by u and draw an edge from u to u' . We also draw edges from u'_1 to u'_2 whenever there is an edge from u_1 to u_2 .

In this manner we obtain all the Hasse diagrams for the unlabeled topologies on the set of $n + 1$ elements.

Definition 8. Let $\mathcal{T}_n = \{T_0 \text{ topologies on } X_n\}$. Let $S_n = \{(T, A) | T \in \mathcal{T}_n \text{ and } A \in T\}$. Define the mapping $L : S_n \rightarrow \mathcal{T}_{n+1}$ by

$$L(T, A) = \{B | B \in T\} \cup \{B \cup \{x_{n+1}\} | B \in T, A \subseteq B\},$$

for some $x_{n+1} \notin X_n$.

Theorem 3. *L is onto.*

Proof. Given a T_0 topology T on a set with $n + 1$ elements we follow the notation of the preceding discussion. We find an element x_{n+1} which is the last to occur as an element of a set of some MC-chain, the T_0 topology T^* on $X_n = X_{n+1} - \{x_{n+1}\}$ and a minimal set $W \in T^*$ such that $W \cup \{x_{n+1}\} \in T$. Then $L(T^*, W) = T$.

□

Then we can recover T from T^* and W by the relation $T = T^* \cup \{A \cup \{x_{n+1}\} | A \in T^* \text{ and } W \subset A\}$.

Thus we have found a recursion such that, given the T_0 topologies on X_n , all of the T_0 topologies on X_{n+1} can be produced. There is necessarily some double counting, as will be shown in the following example.

Example 8. Let $\mathcal{T}_2 = \{T_{2,1}, T_{2,2}, T_{2,3}\}$ be the set of T_0 topologies on the set X_2 , where $T_{2,1} = \{\emptyset, \{a\}, X_2\}$, $T_{2,2} = \{\emptyset, \{b\}, X_2\}$, $T_{2,3} = \{\emptyset, \{a\}, \{b\}, X_2\}$ and $X_2 = \{a, b\}$. Choose $c \notin X_2$. I will show that the recursion defined above produces all five inequivalent T_0 topologies on X_3 .

First we will apply the recursion to $T_{2,1}$. Note that since $T_{2,1}$ is homeomorphic to $T_{2,2}$, the T_0 topologies produced by $T_{2,2}$, in this process would also be homeomorphic to those produced by $T_{2,1}$. For that reason, the process is not shown for $T_{2,2}$. In the Figure 13 below, we selected in turn each element A of $T_{2,1}$. From there we construct the Hasse

diagrams for a T_0 topology on X_3 by letting $S = \{\{c\} \cup B \mid A \subseteq B \in T_{2,1}\}$ and then taking the union of S and $T_{2,1}$.

The process is repeated then repeated for $T_{2,3}$ in Figure 14.

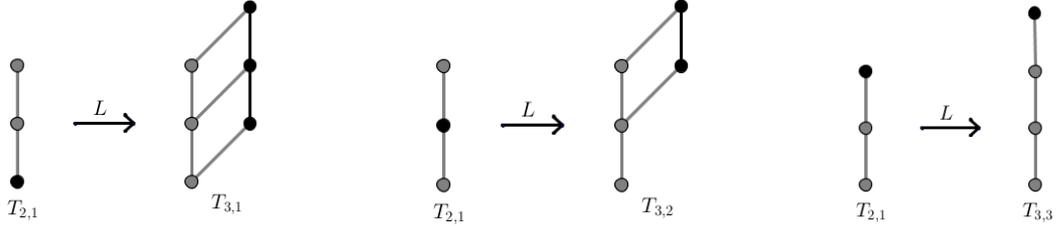


Figure 13: Recursion on $T_{2,1}$ and the produced $T_{3,i}$ topologies, A, S highlighted.

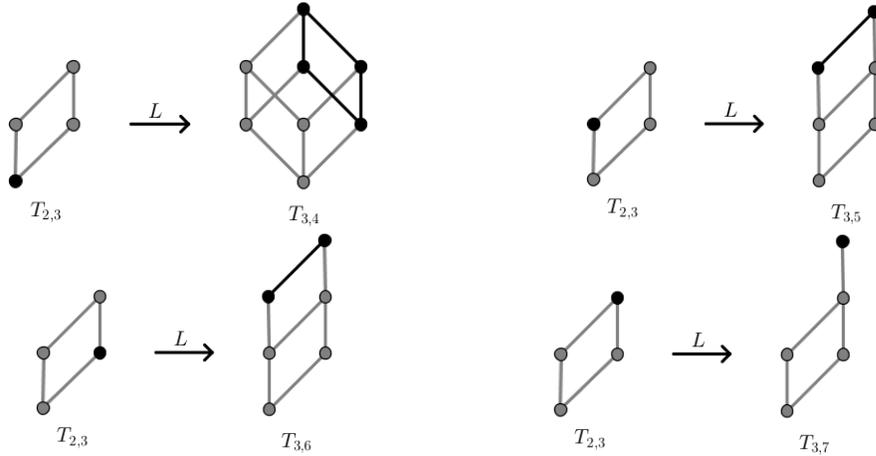


Figure 14: Recursion on $T_{2,2}$ and the produced $T_{3,i}$ topologies, A, S highlighted.

As can be seen, $T_{3,1} \cong T_{3,5} \cong T_{3,6}$, yet all the elements of \mathcal{T}_3 are produced.

We also have from our previous formula which relates the number of topologies on a set X_n to the number of T_0 topologies on $\{X_0, X_1, \dots, X_n\}$ that all non- T_0 topologies are produced from T_0 topologies. It is in this way that the Stirling numbers of the Second Kind $S(n, k)$ are required [7], as $S(n, k)$ is the number of partitions of n elements into k open sets.

Example 9. Consider the indiscrete topology on X_n , $T = \{\emptyset, X_n\}$. When $n = 1$, $T = \{\emptyset, \{a\}\}$ and T is the power set on X_1 , therefore T_0 . It is from this T_0 topology that the

indiscrete topology is produced for all $n > 1$. There is only one way to partition n elements into 1 set, since no element can be placed in $\{\emptyset\}$. Thus, $S(n, 1) = 1$ for all n .

As n increases, the complexity of producing non- T_0 topologies from T_0 topologies increases as well.

Example 10. Consider the T_0 topology $T = \{\emptyset, \{a\}, \{b\}, \{a, b\}\}$ on X_2 . To produce topologies on X_3 , we find all the possible partitions of $\{a, b, c\}$ into two open sets (as no element may be placed in \emptyset and all elements must be placed in X_3). This gives us the $S(3, 2) = 3$ following topologies on X_3 :

$$T_1 = \{\emptyset, \{a\}, \{b, c\}, \{a, b, c\}\}$$

$$T_2 = \{\emptyset, \{b\}, \{a, c\}, \{a, b, c\}\}$$

$$T_3 = \{\emptyset, \{c\}, \{a, b\}, \{a, b, c\}\},$$

which our recursive definition from Chapter 4 failed to produce.

See Appendix A for the Hasse diagrams for the unlabeled topologies on X_2, X_3, X_4 , produced in this fashion, as well as the diagrams for the unlabeled topologies on X_5 [4].

8. SUMMARY

After some unsuccessful efforts we found a recursive definition for generating all T_0 topologies on finite sets. Our initial goal was to obtain a graphical representation of all the topologies for small finite sets.

Mimicking earlier work, the author initially attempted to construct a graph of each T_0 topology on X_3 , with edges defined by $v_1 \leftrightarrow v_2$ if $v_1 \subset v_2$ or $v_2 \subset v_1$, as seen in Figure 15. As before, $P(X_3)$ is not pictured. For higher values of n this construction quickly becomes too cluttered to be of any value. While there are only 19 inequivalent T_0 topologies on a set of three elements, there are 219 on a set of four elements. It is for this reason that this construction was quickly abandoned in favor of a simpler, more compact approach, as the author had done in abandoning the previous shell structures.

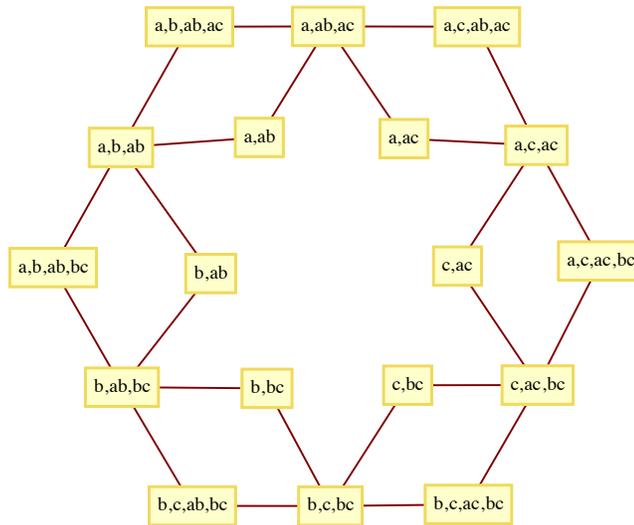


Figure 15: Graph of labeled T_0 topologies on X_3

For that reason, we return to the graph of all MC-chains described in Definition 7, which we denote here by G_n for chains on a base set with n elements.

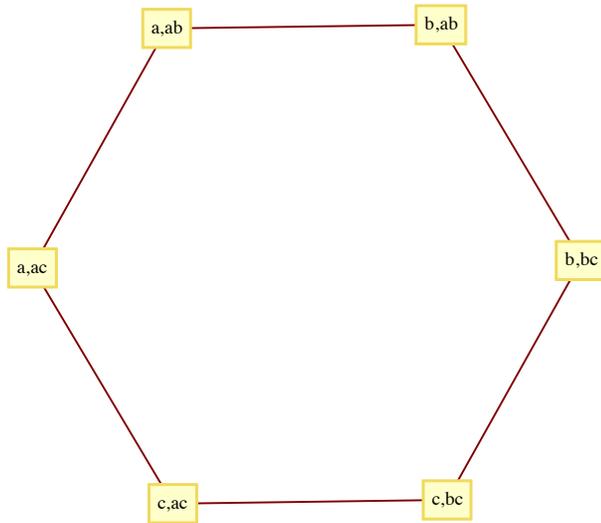


Figure 16: G_3

Example 11. Let $X_3 = \{a, b, c\}$. Then $|V(G_3)| = 6$, and each element of $V(G_3)$ has degree 2. Pictured below in Figure 16 is G_3 .

The construction was then extended to X_4 . The resulting 24-node 3-regular graph is depicted in Figure 17.

While an elegant representation, the graph G_n has some of the same constraints on its analysis as N_n did. Firstly, while every T_0 topology is a connected subgraph of G_n , not every connected subgraph is in turn a T_0 topology. So yet again we encounter the issue of how to specifically construct a subgraph of G_n such that the resulting union of vertices gives a topology. In G_n every vertex represents a minimal T_0 topology, which has $n + 1$ open sets. Every pair of adjacent edges represents a topology with $n + 2$ open sets. Any vertex together with any subset of its immediate neighbors gives a representation of a T_0 topology obtained by taking the union of all the chains represented by those vertices. The simplest connected subsets of G_4 which do not represent topologies are some paths with four vertices. For example, the path $(d, cd, bcd) - (c, cd, bcd) - (c, bc, bcd) - (b, bc, bcd)$ does not represent a topology. The singleton sets with elements b and d are represented,

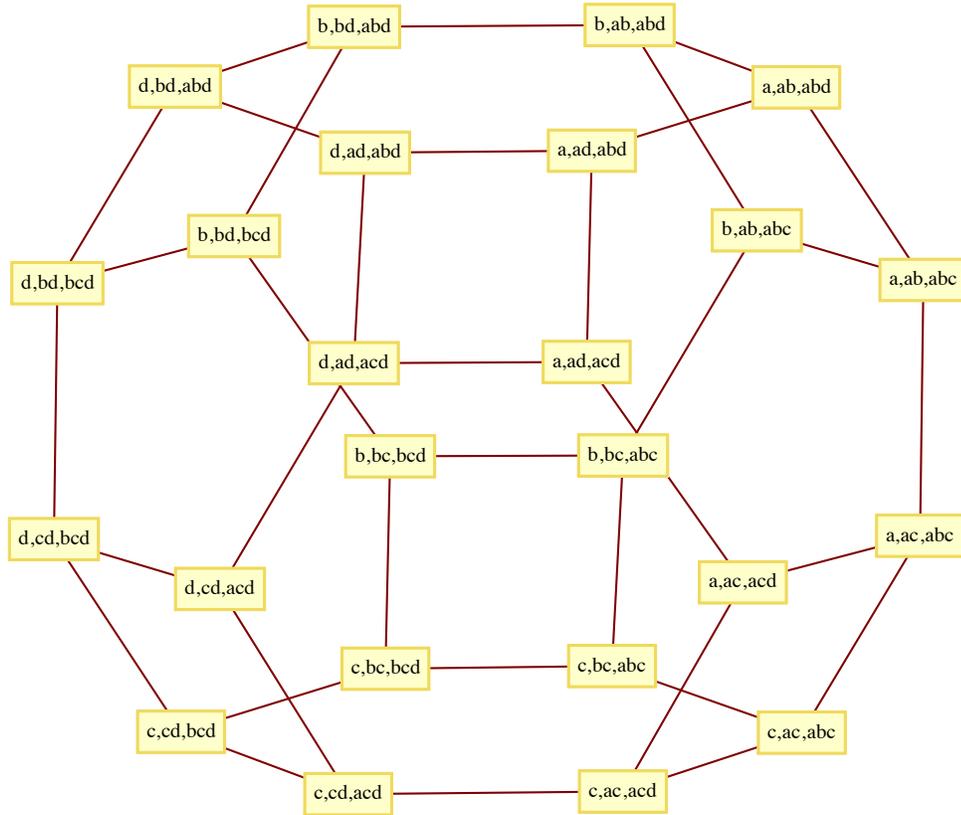


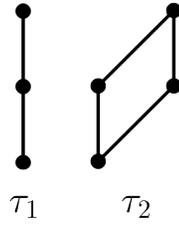
Figure 17: G_4

but not the union of these sets.

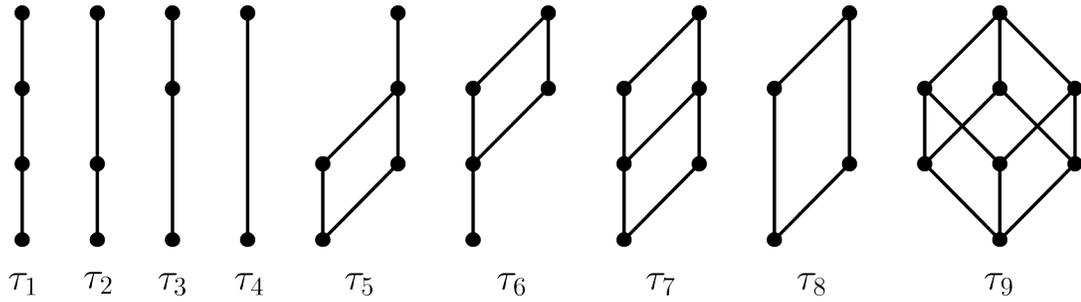
The initial goal of presenting graphical representations of topologies quickly proved intractable. However we did find many interesting relationships between the topologies, and graphical structures which illustrated these relationships for small sets. And while we succeeded in giving a recursive rule to generate all finite T_0 topologies, we can produce a topology on $n + 1$ elements from one on n elements in many different ways. The enumeration of the T_0 topologies remains a challenging problem.

APPENDIX A: UNLABELED HASSE DIAGRAMS OF TOPOLOGIES

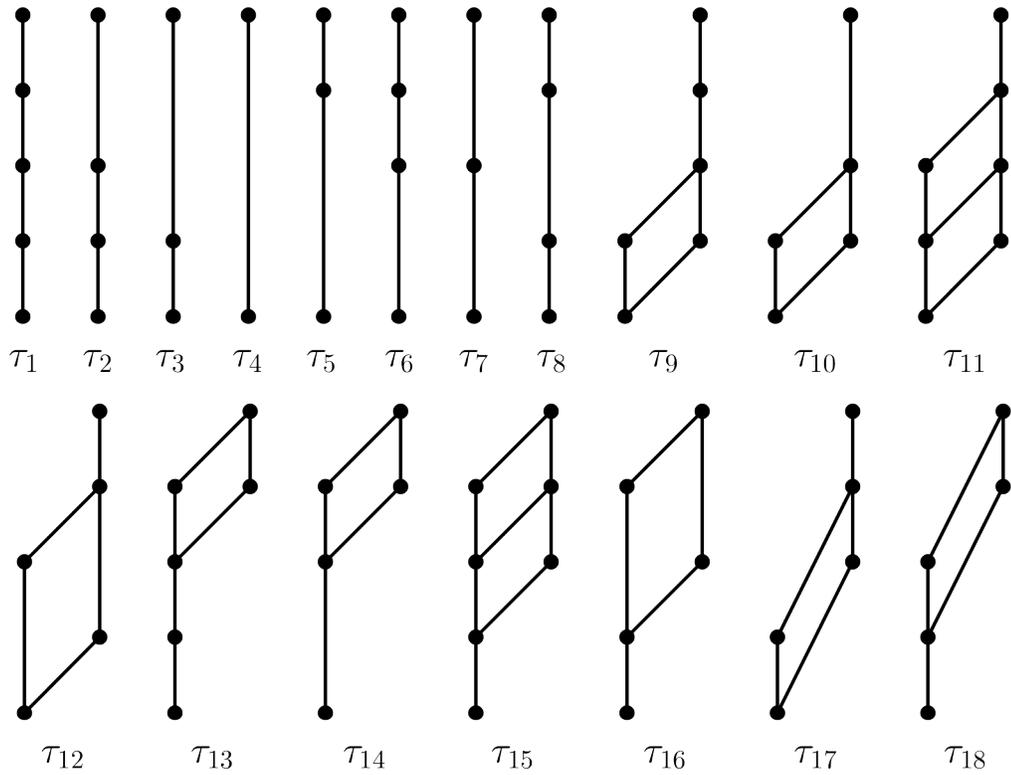
A.1 TOPOLOGIES ON X_2

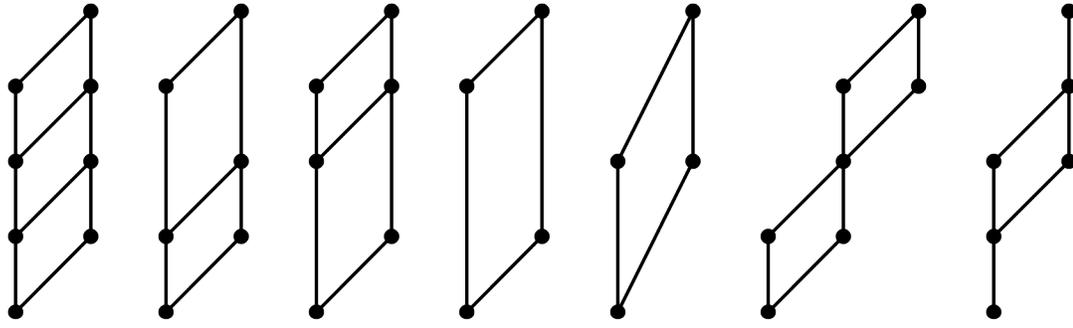


A.2 TOPOLOGIES ON X_3

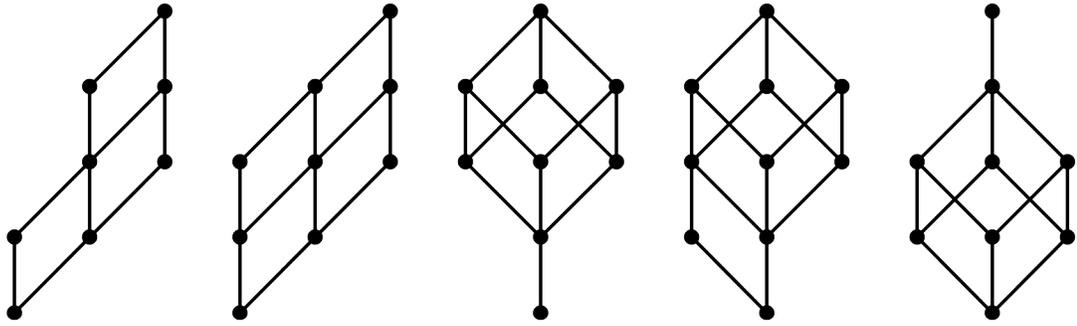


A.3 TOPOLOGIES ON X_4

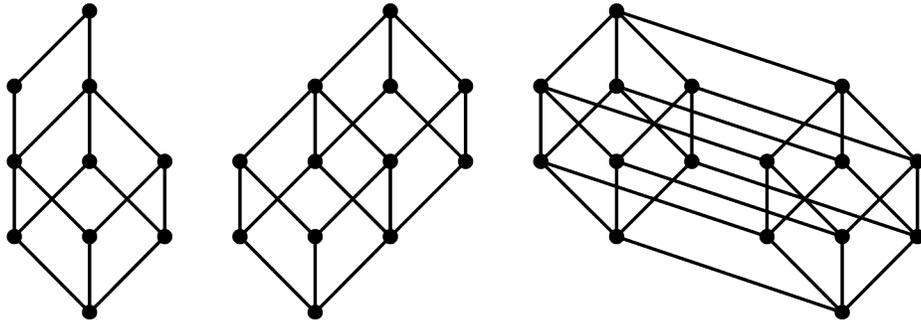




\mathcal{T}_{19} \mathcal{T}_{20} \mathcal{T}_{21} \mathcal{T}_{22} \mathcal{T}_{23} \mathcal{T}_{24} \mathcal{T}_{25}

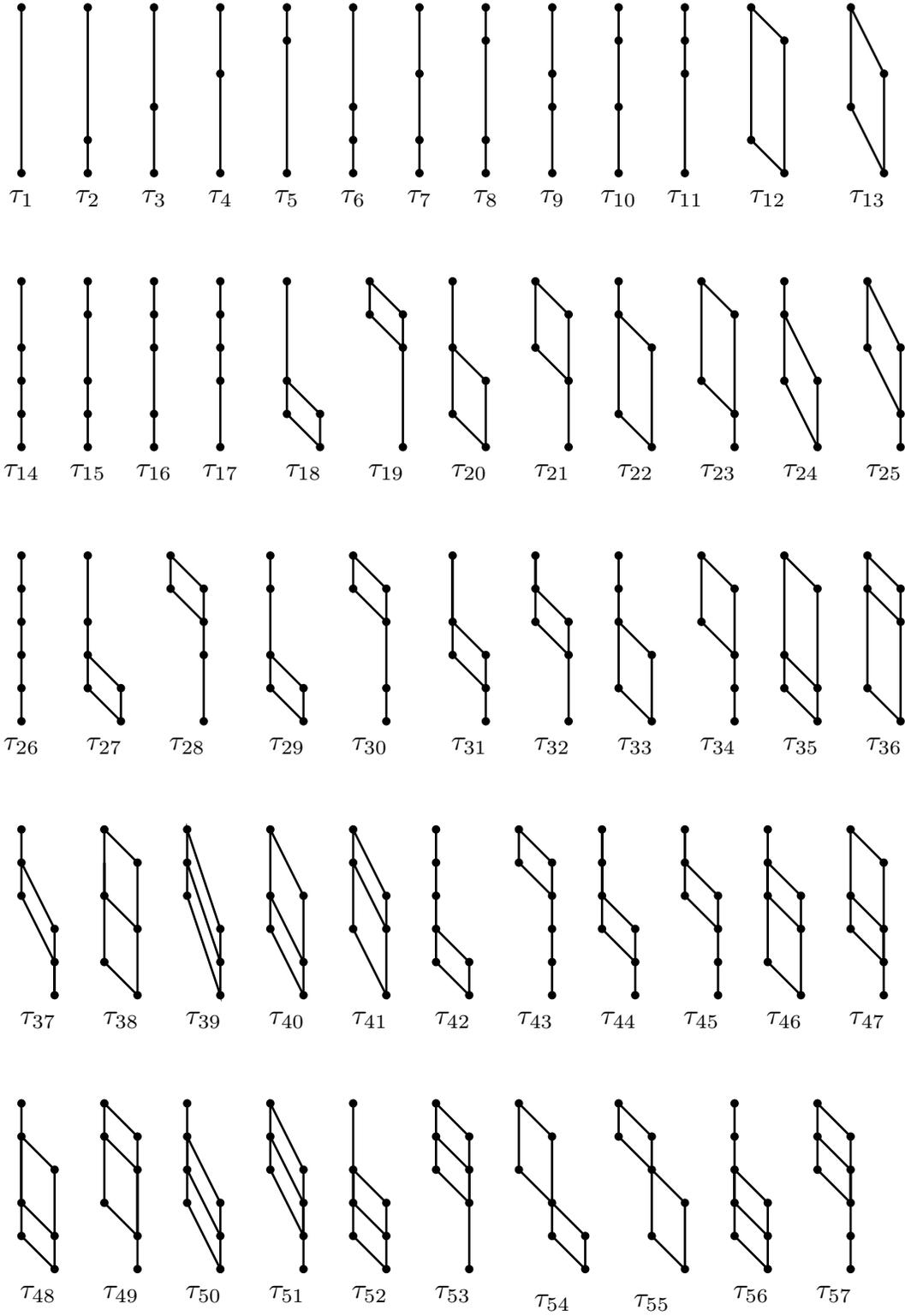


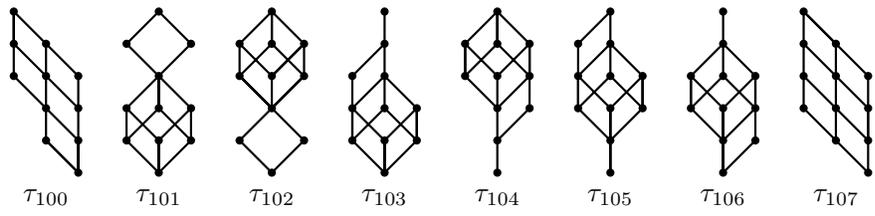
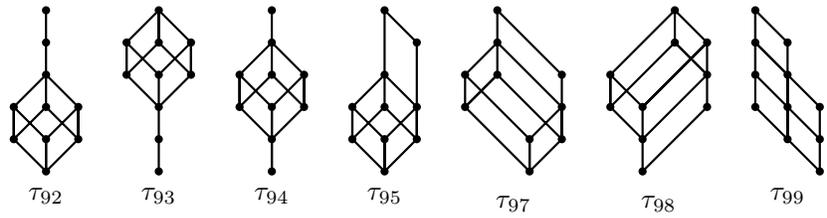
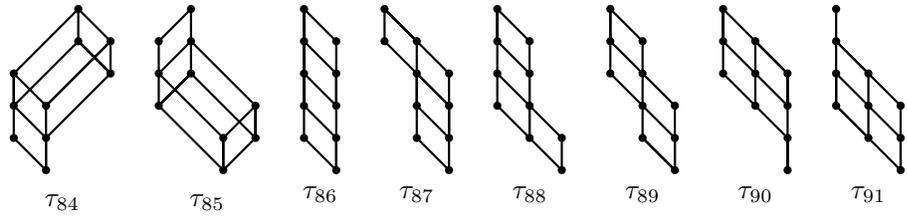
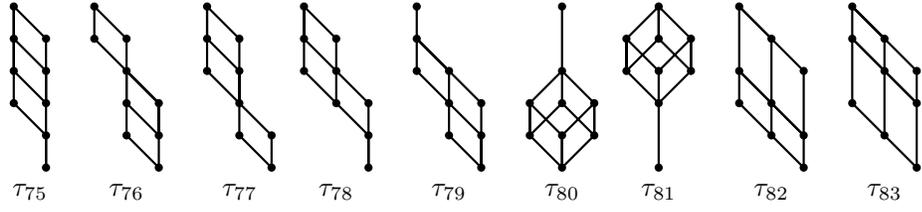
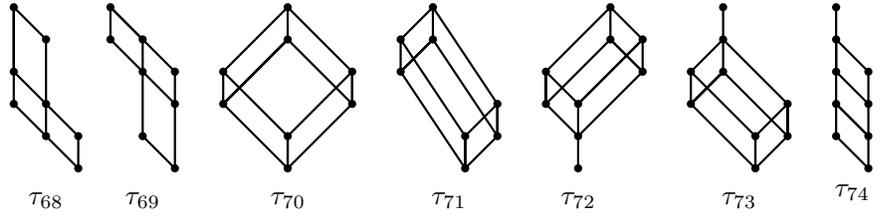
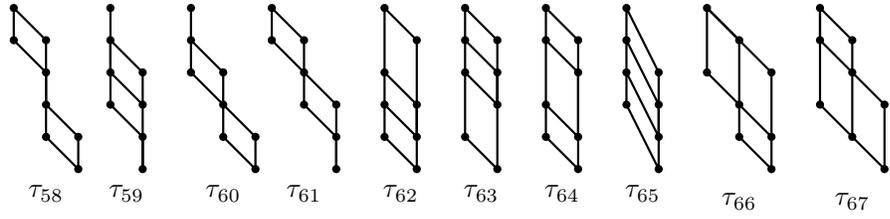
\mathcal{T}_{26} \mathcal{T}_{27} \mathcal{T}_{28} \mathcal{T}_{29} \mathcal{T}_{30}

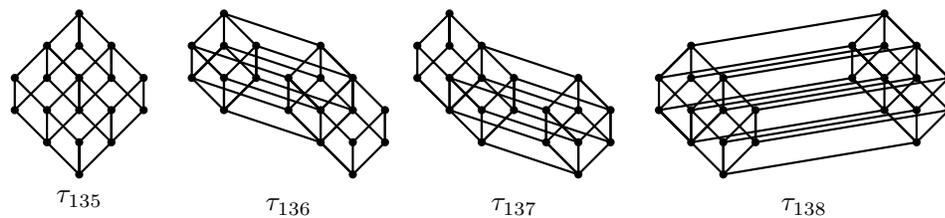
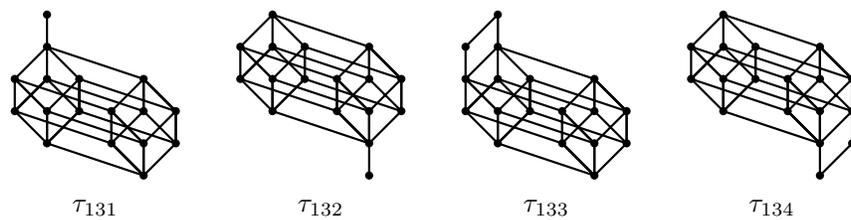
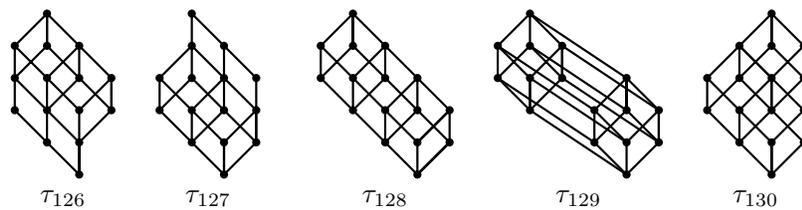
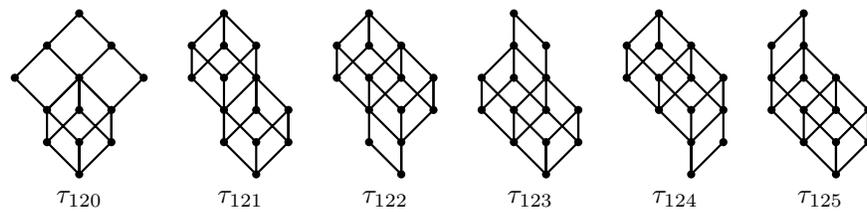
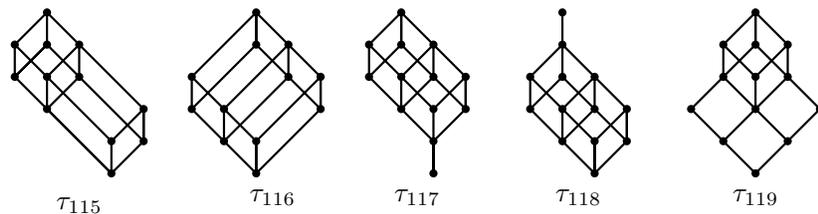
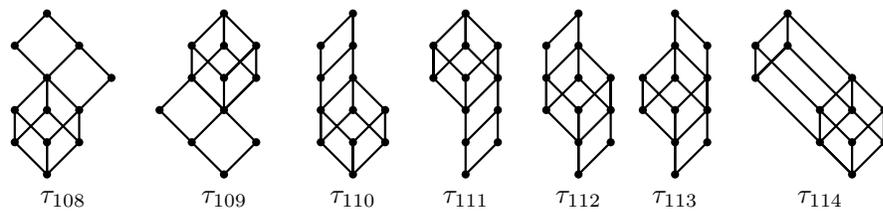


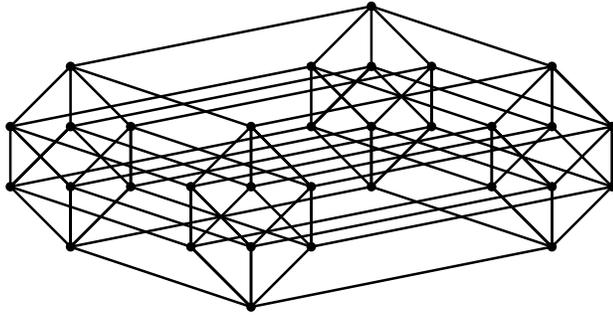
\mathcal{T}_{31} \mathcal{T}_{32} \mathcal{T}_{33}

A.4 TOPOLOGIES ON X_5









τ_{139}

All images for Appendix A.4 were produced by Choo, [4].

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